Evaluation of Relational Operations

Chapter 14
Relational Operations

- We will consider in more detail how to implement:
  - Selection ($\sigma$) Selects a subset of rows from relation.
  - Projection ($\pi$) Deletes unwanted columns from relation.
  - Join ($\bowtie$) Allows us to combine two relations.
  - Set-difference ($-$) Tuples in left reln., but not in right reln.
  - Union ($\cup$) Tuples in reln. 1 and in reln. 2.
  - Aggregation (SUM, MIN, etc.) and GROUP BY

- Since each op returns a relation, ops can be composed!
  After we cover the operations, we will discuss how to optimize queries formed by composing them.
Running Database Example

- Schema

  Sailors (\textit{sid}: integer, \textit{sname}: string, \textit{rating}: integer, \textit{age}: real)
  Reserves (\textit{sid}: integer, \textit{bid}: integer, \textit{day}: dates, \textit{rname}: string)

- \(~100,000\) Reserves:
  - Each tuple is 40 bytes, 100 tuples per page, 1000 pages.

- \(~40,000\) Sailors:
  - Each tuple is 50 bytes, 80 tuples per page, 500 pages.
Selection (review from Chpt 12)

- No Index, Unsorted Data
  - Scan the entire relation, for Reserves $\rightarrow$ 1000 I/Os

- No Index, Sorted Data
  - Binary search, for Reserves $\rightarrow$ $\log_2 1000 \sim 10$ I/Os

- B\textsuperscript{+}-Tree Index, Clustered
  - Use index to find smallest tuple satisfying selection, scan forward from there, for Reserves $\rightarrow$ 3 I/Os to find starting point

- B\textsuperscript{+}-Tree Index, Unclustered
  - Our earlier analyses saw no advantage in this.
Using an Index for Selections

- Cost depends on #qualifying tuples, and clustering.
  - Cost of finding qualifying data entries is typically small, but the cost of retrieving records could be large w/o clustering.
  - Example, assuming uniform distribution of ratings (1-10), about 10% of tuples qualify (100 pages, 10000 tuples). With a clustered index, cost is little more than 100 I/Os; if unclustered, upto 10000 I/Os!

- Important refinement for unclustered indexes:
  1. Find qualifying data entries in index.
  2. Find distinct rids of the pages to be retrieved. (2 ways)
     A. Sort by rid while removing replicates
     B. Build Hash of rids while eliminating replicates
  3. Scan surviving rids while applying selection (result set will be unordered).

- Ensures each page is considered just once (though # of pages is still likely higher than with clustering).
General Selections

- Selections typically involve more than one attribute with logical conjuncts (and, or)
- Recall we transform to CNF (product-of-sum) form
- Can be sorted or clustered by only one attribute
- Only a subset of attributes might have indices
- What order to process selection terms?
- How selective is a selection term?
  - sname = “Joe” < 4% of Sailors
  - age < 20 ~ 10% of Sailors
  - Rating > 7 ~ 30% Sailors
Two Approaches to General Selections

- **First approach**: Find the *most selective access path*, retrieve tuples using it, and apply any remaining terms that don’t *match* the index:
  - *Most selective access path*: An index or file scan that we estimate will require the *fewest page I/Os*.
  - Terms that match this index reduce the number of tuples *retrieved*; other terms are used further discard retrieved tuples, but do not affect number of pages fetched.
  - Consider *day*<8/9/94 AND *bid*=5 AND *sid*=3. A B+ tree index on *day* can be used; then, *bid*=5 and *sid*=3 must be checked for each retrieved tuple. Similarly, a hash index on <*bid*, *sid*> could be used; *day*<8/9/94 must then be checked.
Intersection of Rids

- **Second approach** (if we have 2 or more matching indexes):
  - Get sets of *rids* of data records using each matching index.
  - Then *intersect* these sets of *rids* (we’ll discuss intersection later!)
  - Retrieve the records and apply any remaining terms.
  - Consider *day*<8/9/94 AND *bid*=5 AND *sid*=3. If we have a B+ tree index on *day* and an index on *sid*, both unclustered, we can retrieve distinct *rids* satisfying *day*<8/9/94 using the first, *rids* of recs satisfying *sid*=3 using the second, intersect the *rid sets*, then retrieve records and check *bid*=5.
The Projection Operation

- Modified sorting:
  - Modify Pass 0 of external sort to eliminate unwanted fields. Thus, extending the run-size produced. Tuples in later runs are smaller than input tuples. (Size ratio depends on # and size of fields that are dropped.)
  - Modify merging passes to eliminate duplicates. Thus, number of result tuples smaller than input. (Difference depends on # of duplicates.)
  - Cost: In Pass 0, read original pages, but write out fewer pages (same number of smaller tuples). In merge passes, fewer tuples are written out due to duplicates.

```
SELECT DISTINCT R.sid, R.bid
FROM Reserves R
```
Projection Based on Hashing

- Modified hashing:
  - **Partitioning phase**: Read R using one input buffer. For each tuple, discard unwanted fields, apply hash function $h_1$ to direct output to one of B-1 output buffers.
    - Result is B-1 partitions (of tuples with no unwanted fields). 2 tuples from different partitions guaranteed to be distinct.
  - **Duplicate elimination phase**: For each partition, read it and build an in-memory hash table, using hash function $h_2$ ($\neq h_1$), while discarding duplicates (handled on collisions).
  - **Cost**: For partitioning, read R, write out each tuple, but with fewer fields. This is read in next phase.
Discussion of Projection

- Sort-based approach is the standard; better handling of skewed attribute distributions and result is sorted.
- If an index on the relation contains the wanted projection attributes as its search key, then we use an *index-only* scan (no fetching of the data pages).
- If an ordered (i.e., tree) index contains all wanted attributes in the search key’s *prefix* we can
  - Retrieve data entries in order (index-only scan), discard unwanted fields, compare adjacent tuples to check for duplicates.
Equijoins w/1 common Column

SELECT *  
FROM Reserves R1, Sailors S1  
WHERE R1.sid=S1.sid  

- In algebra: $R \bowtie S$. Very common! Must be carefully optimized. $R \times S$ is large; so, $R \times S$ followed by a selection is inefficient.  
- We will consider more complex join conditions later.  
- Cost metric: # of I/Os. We will ignore output costs.
Simple Nested Loops Join

foreach tuple r in R:
  foreach tuple s in S:
    if r_i == s_j:
      add <r, s> to result

- For each tuple in the outer relation R, we scan the entire inner relation S.
  - Cost: \( M + (p_R \times M) \times N = 1000 + 100 \times 1000 \times 500 \) I/Os.

- Page-at-a-time Nested Loops join: For each page of R, get each page of S, and handle all matching pairs of tuples \(<r, s>\), where r is in R-page and S is in S-page.
  - Cost: \( M + M \times N = 1000 + 1000 \times 500 \)
  - If smaller relation (S) is outer, cost = 500 + 500*1000
Index Nested Loops Join

foreach tuple r in R:
  foreach tuple s in S where r\_i == s\_j:
    add <r, s> to result

- If there is an index on the join column of one relation (say S), make it the inner loop, and exploit the index.
  - Cost: \( M + ( (M*p\_R) * \text{cost of finding matching S tuples}) \)

- For each R tuple, cost of probing S index is about 1.2 for hash index, 2-4 for B+ tree. Cost of then finding S tuples depends on clustering.
  - Clustered index: 1 I/O (typical), unclustered: upto 1 I/O per matching S tuple.
Examples of Index Nested Loops

- Hash-index (Alt. 2) on sid of Sailors (as inner):
  - Scan Reserves: 1000 page I/Os, 100*1000 tuples.
  - For each Reserves tuple: 1.2 I/Os to get data entry in index, plus 1 I/O to get (the exactly one) matching Sailors tuple. Total: 220,000 I/Os.

- Hash-index (Alt. 2) on sid of Reserves (as inner):
  - Scan Sailors: 500 page I/Os, 80*500 tuples.
  - For each Sailors tuple: 1.2 I/Os to find index page with data entries, plus cost of retrieving matching Reserves tuples. Assuming uniform distribution, 2.5 reservations per sailor (100,000 / 40,000). Cost of retrieving them is 1 or 2.5 I/Os depending on whether the index is clustered.
Block Nested Loops Join

- Use one page as an input buffer for scanning the inner S, one page as the output buffer, and use all remaining pages to hold a “block” of outer R.
  - For each matching tuple r in R-block, s in S-page, add <r, s> to result. Then read next R-block, scan S, etc.
Examples of Block Nested Loops

- **Cost:** $M + \left\lfloor M/(B-2) \right\rfloor N$

- With Reserves (R) as outer, and 100 pages of R:
  - Cost of scanning R is 1000 I/Os; a total of 10 blocks.
  - Per block of R, we scan Sailors (S); 10*500 I/Os.
  - If space for just 90 pages of R, we would scan S 12 times.

- With 100-page block of Sailors as outer:
  - Cost of scanning S is 500 I/Os; a total of 5 blocks.
  - Per block of S, we scan Reserves; 5*1000 I/Os.

- Even better to double buffer using a block size of (B-3). Fetching next block while joining current one
Sort-Merge Join \((R \bowtie S)\) (review)

- Sort R and S on the join column, then scan them to “merge” (on join col.), and output result tuples.
  - Advance scan of R until current R-tuple >= current S tuple, then advance scan of S until current S-tuple >= current R tuple; do this until current R tuple = current S tuple.
  - At this point, one-or-more, \(\rho\), R tuples match one-or-more, \(\sigma\), S tuples; output \(<r, s>\) for all pairs of such tuples (\(\rho \times \sigma\)).
  - Then resume scanning R and S.

- Cost: \(M \log M + N \log N + (M+N)\)
Refinements of Sort-Merge Join

- Combine the merging phases of external sorting of R and S with the merging required for the join.
  - Using the sorting refinement that merges multiple runs each pass, we sort R and S up to their last merge pass.
  - Allocate 1 page per run of each relation, and “merge” while checking the join condition.
  - **Cost:** read+writes in (Pass 0.. Pass N-1) + read each relation in (only) merging pass (+ writing of result tuples).
  - Typically reduces I/O cost by a factor of $\frac{1}{2}$.

- In practice, cost of sort-merge join, like the cost of external sorting, is nearly linear.
Hash-Join

- Partition both relations using a common hash function, \( h \), (\( R \) tuples in partition \( i \) will only match \( S \) tuples in partition \( i \)).

- Read in a partition of \( R \), hash it using \( h2 (<> h!) \). Scan matching partition of \( S \), search for matches.
Observations on Hash-Join

- We want each partition of R to fit in B-2 buffer pages, so \(#\text{partitions}, k = \frac{M}{(B - 2)}\), if we assume no skew.
- If we build an in-memory hash table to speed up the matching of tuples, a little more memory is needed.
- If the hash function does not partition uniformly, one or more R partitions may not fit in memory. Can apply hash-join technique recursively to this partition and do the join of this R-partition with corresponding S-partition.
Cost of Hash-Join

- In partitioning phase, read+write both relns: $2(M+N)$. In matching phase, read both relns: $M+N$ I/Os.
- In our running example, this is a total of 4500 I/Os.
- Sort-Merge Join vs. Hash Join:
  - Both have a cost of $3(M+N)$ I/Os. Hash-Join is superior if relation sizes differ greatly. Also, Hash-Join shown to be highly parallelizable.
  - Sort-Merge insensitive to data skew; and result is sorted.
General Join Conditions

- **Equalities over several attributes (e.g., $R\text{.}sid=S\text{.}sid$ AND $R\text{.}rname=S\text{.}sname$):**
  - For Index NL, build index on $<\text{sid, sname}>$ (if S is inner); or use existing indexes on $\text{sid}$ or $\text{sname}$.
  - For Sort-Merge and Hash Join, sort/partition on combination of the two join columns.

- **Inequality conditions (e.g., $R\text{.}rname < S\text{.}sname$):**
  - For Index NL, need (clustered!) B+ tree index.
    - Perform range probes on inner; # matches likely to be much higher than for equality joins.
  - Hash Join, Sort Merge Join not applicable.
  - Block NL quite likely to be the best join method here.
Set Operations

- Intersection and cross-product special cases of join.
- Union (Distinct) and Except similar; we’ll do union.
- Sorting based approach to union:
  - Sort both relations (on combination of all attributes).
  - Scan sorted relations and merge them.
  - Alternative: Merge runs from Pass 0 for both relations.
- Hash based approach to union:
  - Partition R and S using hash function $h$.
  - For each S-partition, build in-memory hash table (using $h2$), scan corr. R-partition and add tuples to table while discarding duplicates.
Aggregate Operations (AVG, MIN, etc.)

- Without grouping:
  - In general, requires scanning the relation.
  - Given index whose search key includes all attributes in the SELECT or WHERE clauses, can do index-only scan.

- With grouping:
  - Sort on group-by attributes, then scan relation and compute aggregate for each group. (Can improve upon this by combining sorting and aggregate computation.)
  - Similar approach based on hashing on group-by attributes.
  - Given tree index whose search key includes all attributes in SELECT, WHERE and GROUP BY clauses, can do index-only scan; if group-by attributes form prefix of search key, can retrieve data entries/tuples in group-by order.
Impact of Buffering

- If several operations are executing concurrently, estimating the number of available buffer pages is guesswork.
- Repeated access patterns interact with buffer replacement policy.
  - e.g., Inner relation is scanned repeatedly in Simple Nested Loop Join. With enough buffer pages to hold inner, replacement policy does not matter. Otherwise, MRU is best, LRU is worst (sequential flooding).
  - Does replacement policy matter for Block Nested Loops?
  - What about Index Nested Loops? Sort-Merge Join?
Summary

- A virtue of relational DBMSs: *queries are composed of a few basic operators*; the implementation of these operators can be carefully tuned (and it is important to do this!).

- Many alternative implementation techniques for each operator; no universally superior technique for most operators.

- Must consider available alternatives for each operation in a query and choose best one based on system statistics, etc. This is part of the broader task of optimizing a query composed of several ops.